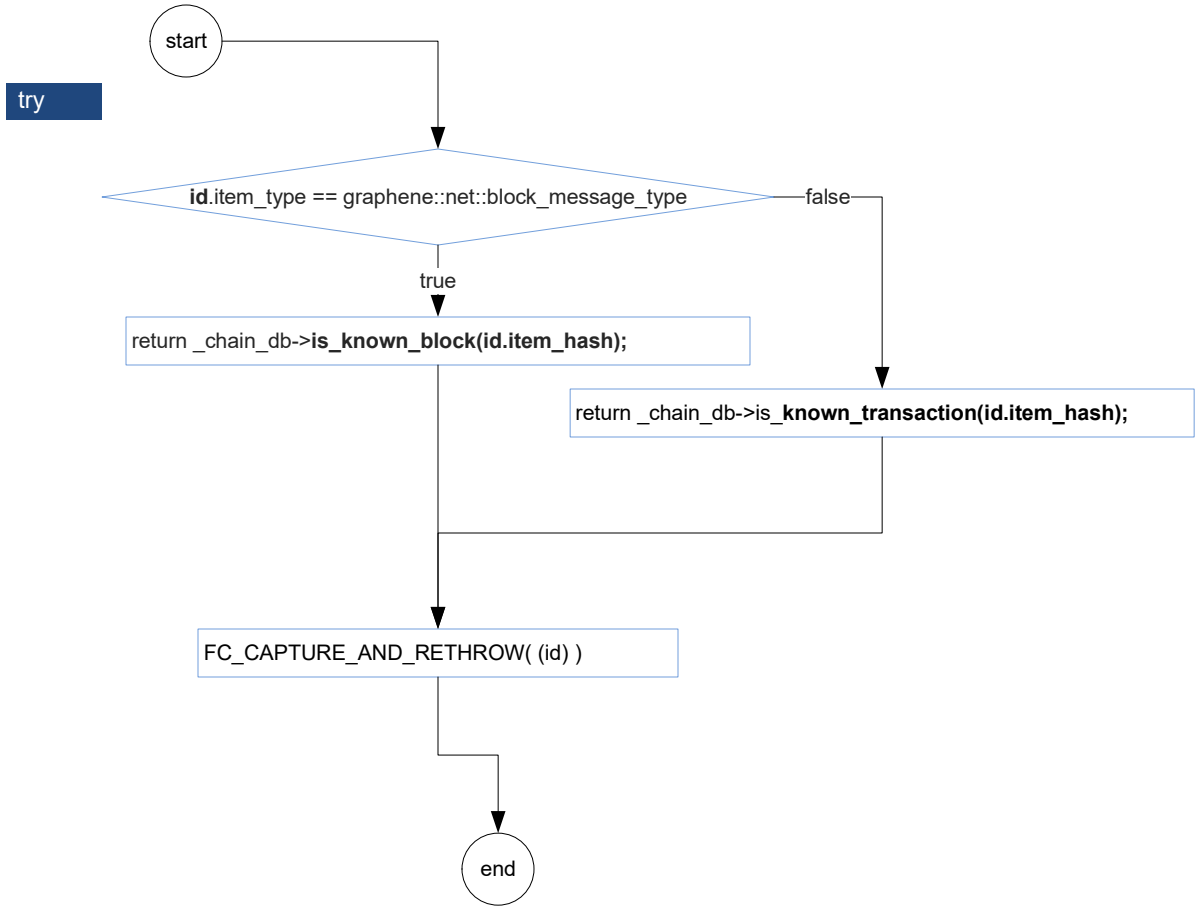


* If delegate has the item, the network has no need to fetch it.



```
bool application_impl::handle_block(const graphene::net::block_message& blk_msg,
                                   bool sync_mode,
                                   std::vector<fc::uint160_t>& contained_transaction_message_ids)
```

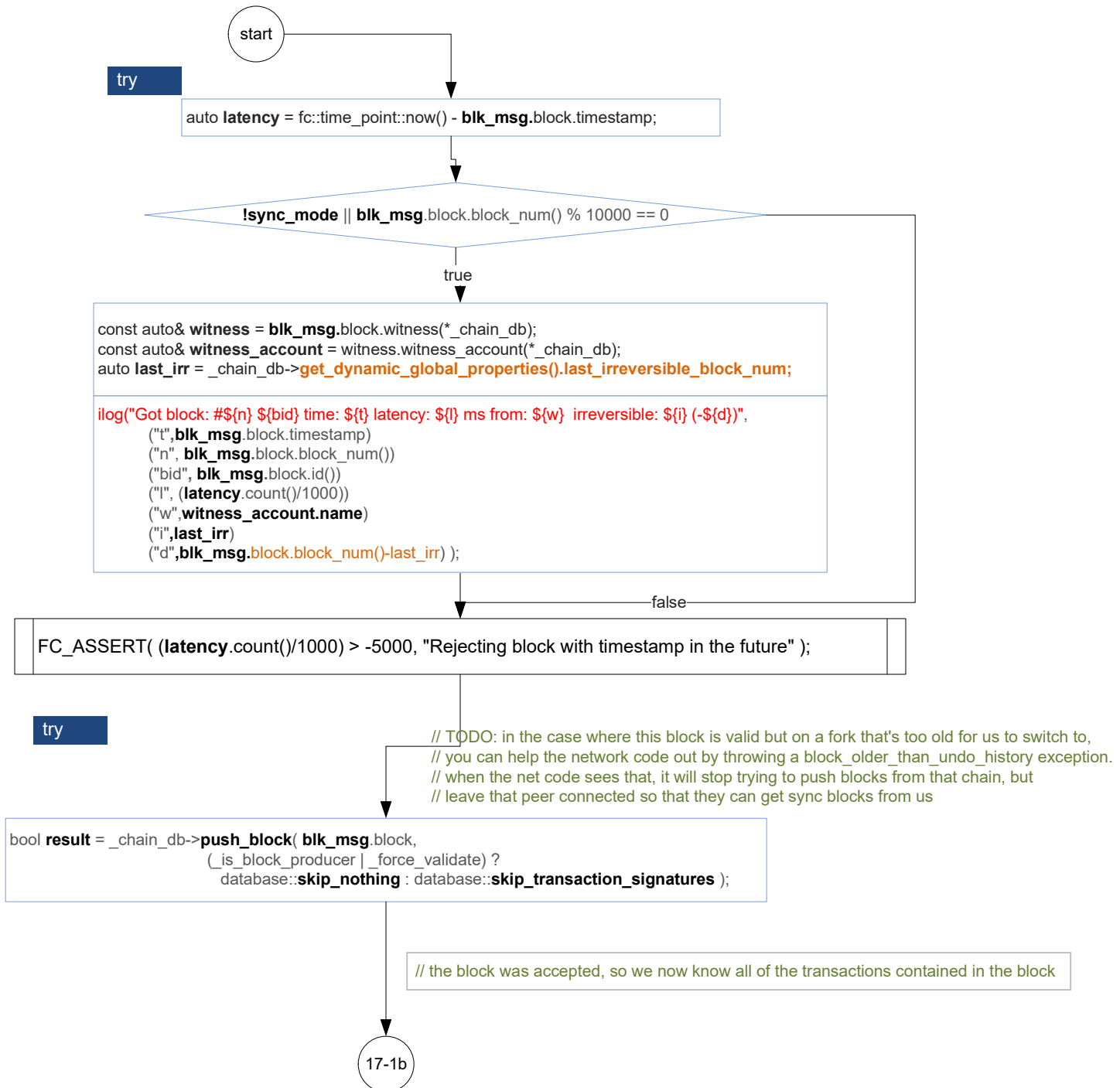
application.cpp (17-1)

@brief allows the application to validate an item prior to broadcasting to peers.

@param sync_mode true if the message was fetched through the sync process, false during normal operation

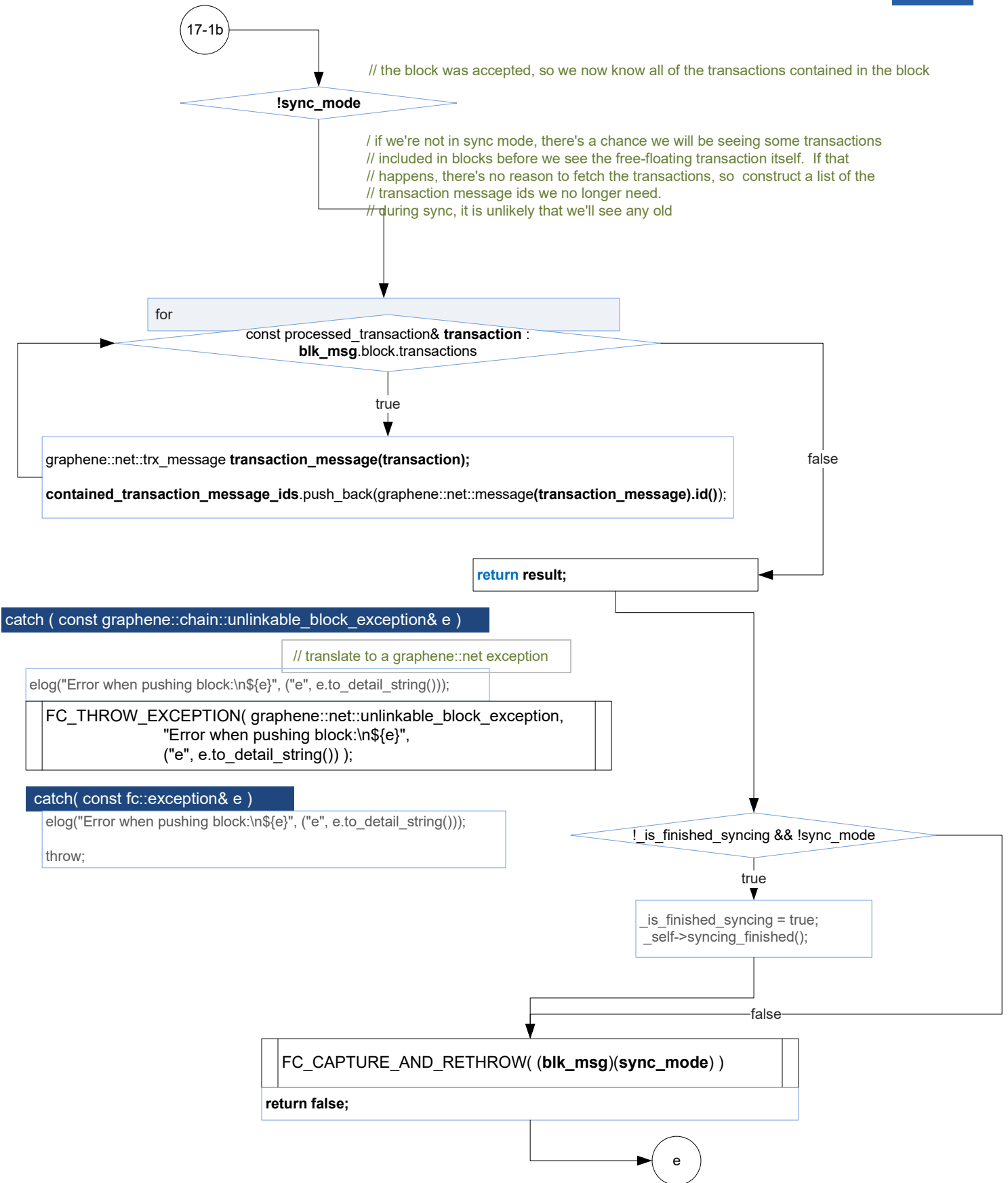
@returns true if this message caused the blockchain to switch forks, false if it did not

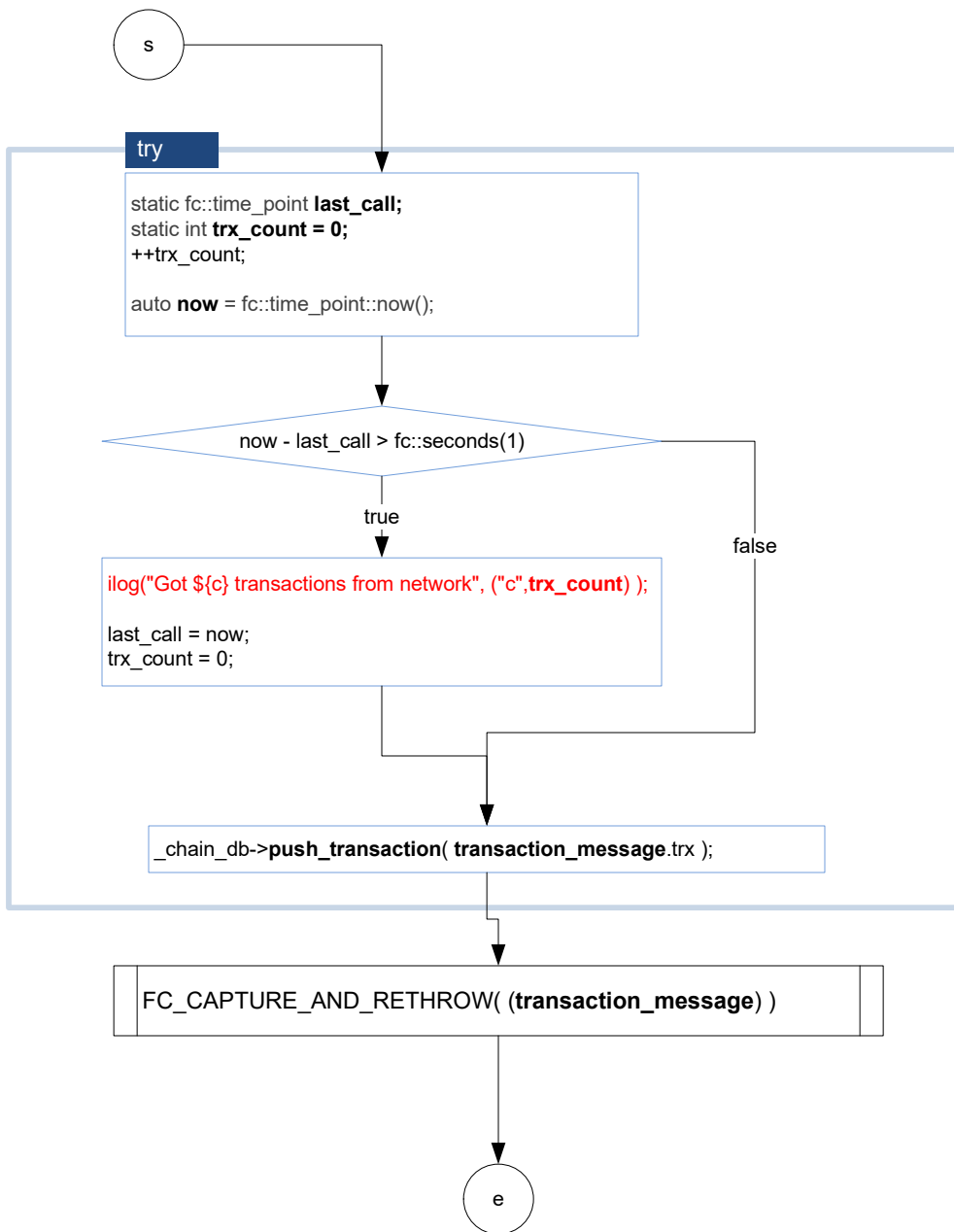
@throws exception if error validating the item, otherwise the item is safe to broadcast on.

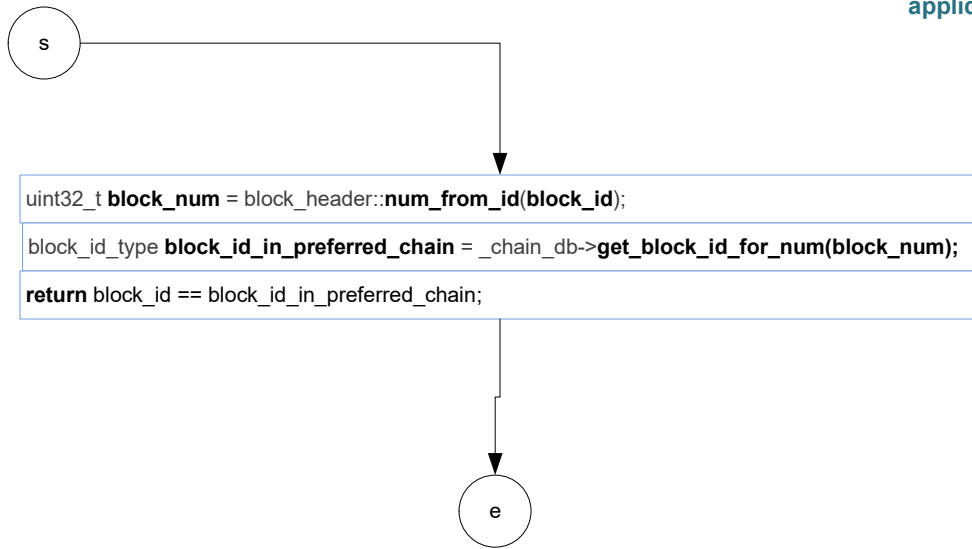


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bool application_impl::handle_block(const graphene::net::block_message& blk_msg,
    bool sync_mode,
    std::vector<fc::uint160_t>& contained_transaction_message_ids)
```

application.cpp (17-2)

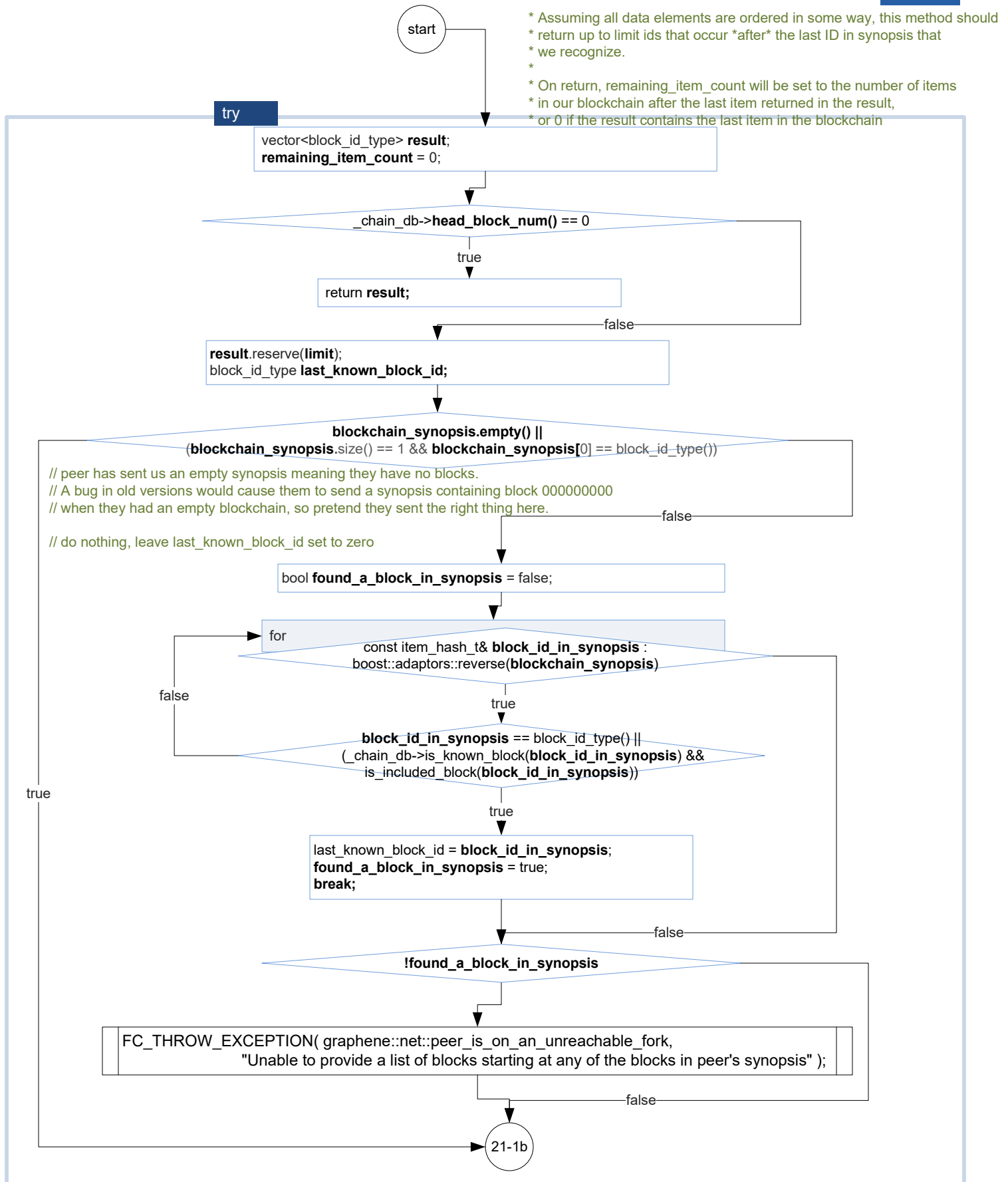






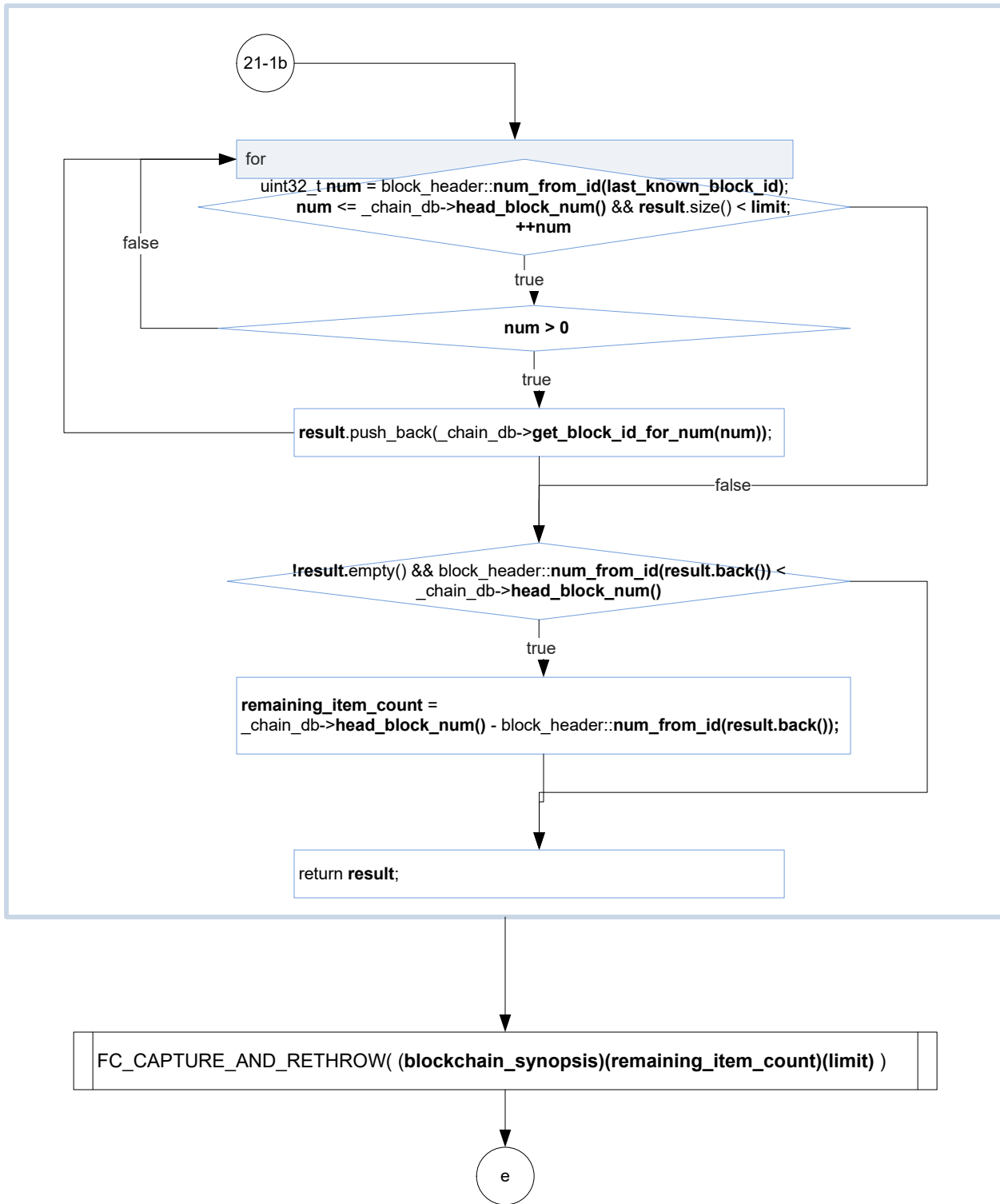
```
std::vector<item_hash_t> application_impl::get_block_ids(const std::vector<item_hash_t>& blockchain_synopsis,
uint32_t& remaining_item_count,
uint32_t limit)
```

application.cpp (21-1)



```
std::vector<item_hash_t> application_impl::get_block_ids(const std::vector<item_hash_t>& blockchain_synopsis,  
uint32_t& remaining_item_count,  
uint32_t limit)
```

application.cpp (21-2)




```
std::vector<item_hash_t> application_impl::get_blockchain_synopsis(const item_hash_t& reference_point,
    uint32_t number_of_blocks_after_reference_point)
```

application.cpp

```
/**
 * Returns a synopsis of the blockchain used for syncing. This consists of a list of
 * block hashes at intervals exponentially increasing towards the genesis block.
 * When syncing to a peer, the peer uses this data to determine if we're on the same
 * fork as they are, and if not, what blocks they need to send us to get us on their
 * fork.
 *
 * In the over-simplified case, this is a straightforward synopsis of our current
 * preferred blockchain; when we first connect up to a peer, this is what we will be sending.
 * It looks like this:
 * If the blockchain is empty, it will return the empty list.
 * If the blockchain has one block, it will return a list containing just that block.
 * If it contains more than one block:
 * the first element in the list will be the hash of the highest numbered block that
 * we cannot undo
 * the second element will be the hash of an item at the half way point in the undoable
 * segment of the blockchain
 * the third will be ~3/4 of the way through the undoable segment of the block chain
 * the fourth will be at ~7/8...
 * &c.
 * the last item in the list will be the hash of the most recent block on our preferred chain
 * so if the blockchain had 26 blocks labeled a - z, the synopsis would be:
 * a n u x z
 * the idea being that by sending a small (<30) number of block ids, we can summarize a huge
 * blockchain. The block ids are more dense near the end of the chain where because we are
 * more likely to be almost in sync when we first connect, and forks are likely to be short.
 * If the peer we're syncing with in our example is on a fork that started at block 'v',
 * then they will reply to our synopsis with a list of all blocks starting from block 'u',
 * the last block they know that we had in common.
 *
 * In the real code, there are several complications.
 *
 * First, as an optimization, we don't usually send a synopsis of the entire blockchain, we
 * send a synopsis of only the segment of the blockchain that we have undo data for. If their
 * fork doesn't build off of something in our undo history, we would be unable to switch, so there's
 * no reason to fetch the blocks.
 *
 * Second, when a peer replies to our initial synopsis and gives us a list of the blocks they think
 * we are missing, they only send a chunk of a few thousand blocks at once. After we get those
 * block ids, we need to request more blocks by sending another synopsis (we can't just say "send me
 * the next 2000 ids" because they may have switched forks themselves and they don't track what
 * they've sent us). For faster performance, we want to get a fairly long list of block ids first,
 * then start downloading the blocks.
 *
 * The peer doesn't handle these follow-up block id requests any different from the initial request;
 * it treats the synopsis we send as our blockchain and bases its response entirely off that. So to
 * get the response we want (the next chunk of block ids following the last one they sent us, or,
 * failing that, the shortest fork off of the last list of block ids they sent), we need to construct
 * a synopsis as if our blockchain was made up of:
 * 1. the blocks in our block chain up to the fork point (if there is a fork) or the head block (if no fork)
 * 2. the blocks we've already pushed from their fork (if there's a fork)
 * 3. the block ids they've previously sent us
 * Segment 3 is handled in the p2p code, it just tells us the number of blocks it has (in
 * number_of_blocks_after_reference_point) so we can leave space in the synopsis for them.
 *
 * We're responsible for constructing the synopsis of Segments 1 and 2 from our active blockchain and
 * fork database. The reference_point parameter is the last block from that peer that has been
 * successfully pushed to the blockchain, so that tells us whether the peer is on a fork or on
 * the main chain.
 */
```

```
std::vector<item_hash_t> application_impl::get_blockchain_synopsis(const item_hash_t& reference_point,
    uint32_t number_of_blocks_after_reference_point)
```



try

```
std::vector<item_hash_t> synopsis;
synopsis.reserve(30);
uint32_t high_block_num;
uint32_t non_fork_high_block_num;
uint32_t low_block_num = _chain_db->last_non_undoable_block_num();
std::vector<block_id_type> fork_history;
```

reference_point != item_hash_t()

false

*// the node is asking for a summary of the block chain up to a specified
 // block, which may or may not be on a fork
 // for now, assume it's not on a fork*



is_included_block(reference_point)

false

// reference_point is a block we know about and is on the main chain



```
uint32_t reference_point_block_num = block_header::num_from_id(reference_point);
assert(reference_point_block_num > 0);

high_block_num = reference_point_block_num;
non_fork_high_block_num = high_block_num;
```

reference_point_block_num < low_block_num

false



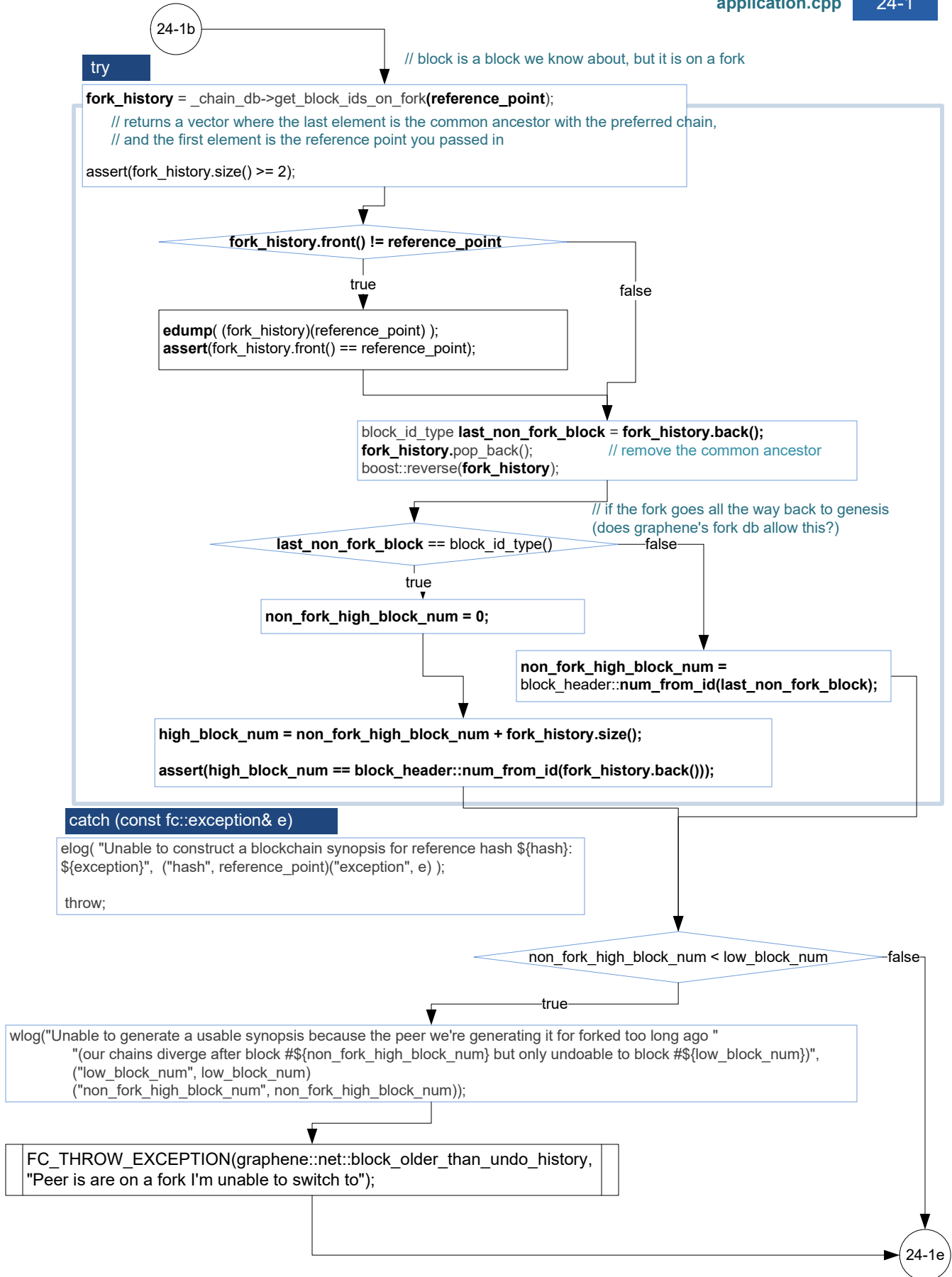
*// we're on the same fork (at least as far as reference_point) but we've passed
 // reference point and could no longer undo that far if we diverged after that
 // block. This should probably only happen due to a race condition where
 // the network thread calls this function, and then immediately pushes a bunch of blocks,
 // then the main thread finally processes this function.
 // with the current framework, there's not much we can do to tell the network
 // thread what our current head block is, so we'll just pretend that
 // our head is actually the reference point.
 // this *may* enable us to fetch blocks that we're unable to push, but that should
 // be a rare case (and correctly handled)*

```
low_block_num = reference_point_block_num;
```

true



```
std::vector<item_hash_t> application_impl::get_blockchain_synopsis(const item_hash_t& reference_point,
    uint32_t number_of_blocks_after_reference_point)
```



```
std::vector<item_hash_t> application_impl::get_blockchain_synopsis(const item_hash_t& reference_point,
    uint32_t number_of_blocks_after_reference_point)
```

